# CS 70Discrete Mathematics and Probability TheoryFall 2019Alistair Sinclair and Yun S. SongDIS 4

## 1 Modular Practice

Solve the following modular arithmetic equations for *x* and *y*.

- (a)  $9x + 5 \equiv 7 \pmod{11}$ .
- (b) Show that  $3x + 15 \equiv 4 \pmod{21}$  does not have a solution.
- (c) The system of simultaneous equations  $3x + 2y \equiv 0 \pmod{7}$  and  $2x + y \equiv 4 \pmod{7}$ .

- (d)  $13^{2019} \equiv x \pmod{12}$ .
- (e)  $7^{67} \equiv x \pmod{11}$ .

#### 2 Fibonacci GCD

The Fibonacci sequence is given by  $F_n = F_{n-1} + F_{n-2}$ , where  $F_0 = 0$  and  $F_1 = 1$ . Prove that, for all  $n \ge 0$ ,  $gcd(F_n, F_{n-1}) = 1$ .

## 3 RSA Warm-Up

Consider an RSA scheme with modulus N = pq, where p and q are distinct prime numbers larger than 3.

- (a) What is wrong with using the exponent e = 2 in an RSA public key?
- (b) Recall that *e* must be relatively prime to p-1 and q-1. Find a condition on *p* and *q* such that e = 3 is a valid exponent.
- (c) Now suppose that p = 5, q = 17, and e = 3. What is the public key?
- (d) What is the private key?
- (e) Alice wants to send a message x = 10 to Bob. What is the encrypted message E(x) she sends using the public key?
- (f) Suppose Bob receives the message y = 24 from Alice. What equation would he use to decrypt the message? and what is the decrypted message?

#### 4 Breaking RSA

Eve is not convinced she needs to factor N = pq in order to break RSA. She argues: "All I need to know is (p-1)(q-1)... then I can find *d* as the inverse of  $e \mod (p-1)(q-1)$ . This should be easier than factoring *N*." Prove Eve wrong, by showing that if she knows (p-1)(q-1), she can easily factor *N* (thus showing finding (p-1)(q-1) is at least as hard as factoring *N*).